# ECE544: Fault-Tolerant Computing \& Reliability Engineering 

 (Fall 2022)
## Homework \#2 Solution

 (45 points)1. How many check bits are needed if the Hamming correcting code is used to detect single bit errors in a 64-bit data word? ( 5 points)

## Solution:

Need $K$ check bits such that: $64+K \leq 2^{K}-1$.
The minimum value of K , which satisfies this condition, is 7 .
2. Develop an SEC code for a 16-bit data word. 1) Generate the code for the data word 0101000000111001 . 2) Show that the code will correctly identify an error in data bit $\mathrm{D}_{16}$. (40 points)

## Solution:

1) (30 points)

Step 1 ( 4 points): According to the inequality $2^{k}-1>=M+K$, where $M=16,5$ check bits are needed for an SEC code for 16-bit data words.

Step 2 ( 9 points): The layout of data bits and check bits:

| Bit Position | Position Number | Check Bits | Data Bits |
| :---: | :---: | :---: | :---: |
| 21 | 10101 |  | $\mathrm{D}_{16}$ |
| 20 | 10100 |  | $\mathrm{D}_{15}$ |
| 19 | 10011 |  | $\mathrm{D}_{14}$ |
| 18 | 10010 |  | $\mathrm{D}_{13}$ |
| 17 | 10001 | C 16 | $\mathrm{D}_{12}$ |
| 16 | 10000 |  | $\mathrm{D}_{11}$ |
| 15 | 01111 |  | $\mathrm{D}_{10}$ |
| 14 | 01110 |  | $\mathrm{D}_{9}$ |
| 13 | 01101 |  | $\mathrm{D}_{8}$ |
| 12 | 01100 |  | $\mathrm{D}_{7}$ |
| 11 | 01011 |  | $\mathrm{D}_{6}$ |
| 10 | 01010 |  | $\mathrm{D}_{5}$ |
| 9 | 01001 | C 8 |  |
| 8 | 01000 |  | $\mathrm{D}_{4}$ |
| 7 | 00111 |  | $\mathrm{D}_{3}$ |
| 6 | 00110 |  |  |


| 5 | 00101 |  | $\mathrm{D}_{2}$ |
| :--- | :--- | :--- | :--- |
| 4 | 00100 | C 4 |  |
| 3 | 00011 | C 2 | $\mathrm{D}_{1}$ |
| 2 | 00010 | C 1 |  |
| 1 | 00001 |  |  |

Step 3 (15 points): The check bits are calculated
$\mathrm{C} 1=\mathrm{D} 1 \oplus \mathrm{D} 2 \oplus \mathrm{D} 4 \oplus \mathrm{D} 5 \oplus \mathrm{D} 7 \oplus \mathrm{D} 9 \oplus \mathrm{D} 11 \oplus \mathrm{D} 12 \oplus \mathrm{D} 14 \oplus \mathrm{D} 16$
$\mathrm{C} 2=\mathrm{D} 1 \oplus \mathrm{D} 3 \oplus \mathrm{D} 4 \oplus \mathrm{D} 6 \oplus \mathrm{D} 7 \oplus \mathrm{D} 10 \oplus \mathrm{D} 11 \oplus \mathrm{D} 13 \oplus \mathrm{D} 14$
$\mathrm{C} 4=\mathrm{D} 2 \oplus \mathrm{D} 3 \oplus \mathrm{D} 4 \oplus \mathrm{D} 8 \oplus \mathrm{D} 9 \oplus \mathrm{D} 10 \oplus \mathrm{D} 11 \oplus \mathrm{D} 15 \oplus \mathrm{D} 16$
$\mathrm{C} 8=\mathrm{D} 5 \oplus \mathrm{D} 6 \oplus \mathrm{D} 7 \oplus \mathrm{D} 8 \oplus \mathrm{D} 9 \oplus \mathrm{D} 10 \oplus \mathrm{D} 11$
$\mathrm{C} 16=\mathrm{D} 12 \oplus \mathrm{D} 13 \oplus \mathrm{D} 14 \oplus \mathrm{D} 15 \oplus \mathrm{D} 16$
For the word $\mathrm{D}_{16} \mathrm{D}_{15} \mathrm{D}_{14} \ldots \mathrm{D}_{2} \mathrm{D}_{1}=\mathbf{0 1 0 1 0 0 0 0 0 0 1 1 1 0 0 1}$, the check bits are $\mathrm{C} 16=0 ; \mathrm{C} 8=0 ; \mathrm{C} 4=0 ; \mathrm{C} 2=0 ; \mathrm{C} 1=1$.

Step 4 ( 2 points): The code word is $\mathbf{0 1 0 1 0} \mathbf{0 0 0 0 0 1 1 0 1 0 0 0 1 0 1}$

## 2) (10 points)

If an error occurs in data bit $\mathrm{D}_{16}$, the check bits become

$$
\mathrm{C} 16=1 ; \mathrm{C} 8=0 ; \mathrm{C} 4=1 ; \mathrm{C} 2=0 ; \mathrm{C} 1=0 .
$$

Comparing the two sets of check bits forms the syndrome word:


The result indicates an error identified in bit position 21, which is data bit $\mathrm{D}_{16}$

